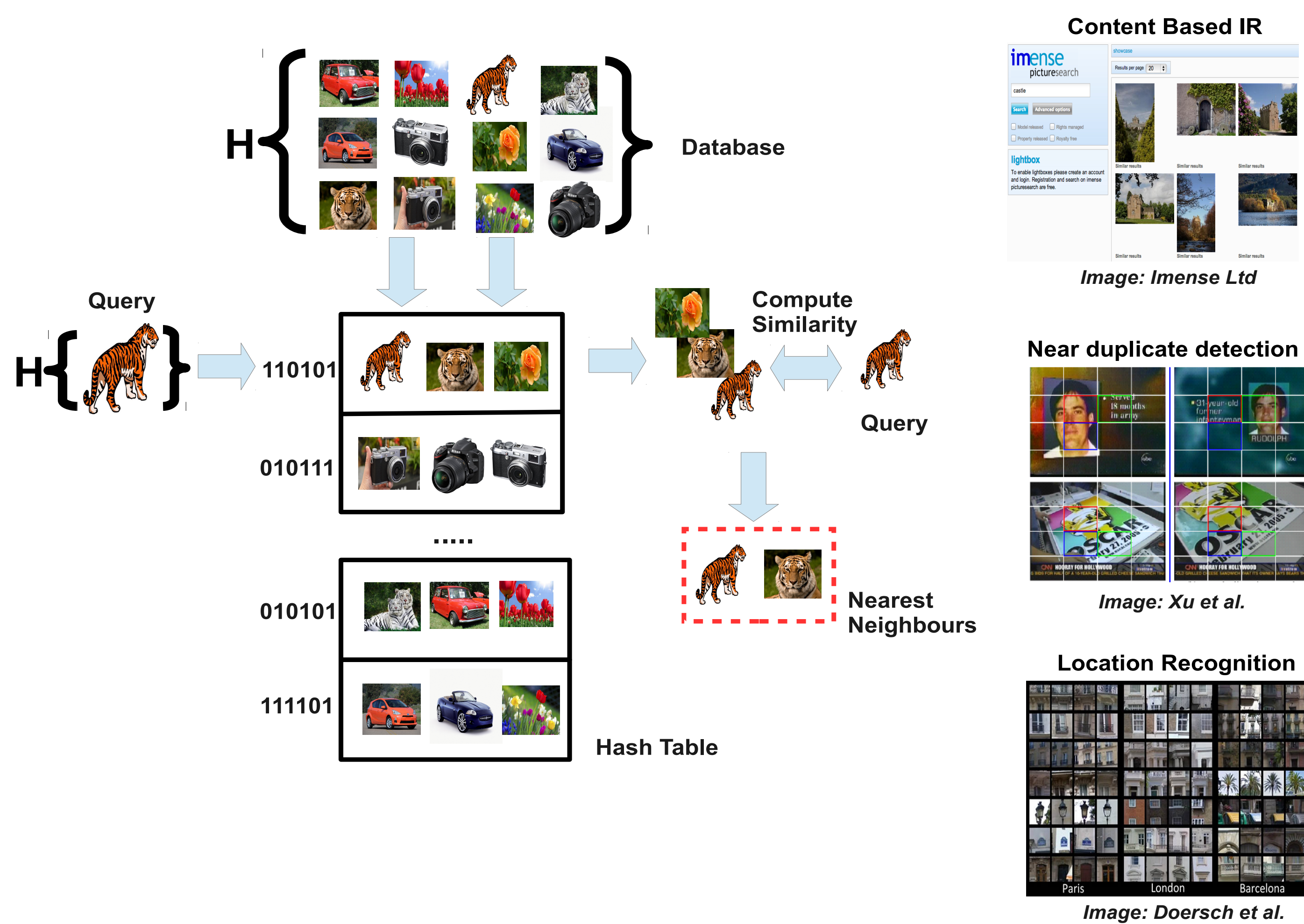


RESEARCH QUESTION

- LSH uses 1 bit per hyperplane. Can we do better with multiple bits?

INTRODUCTION

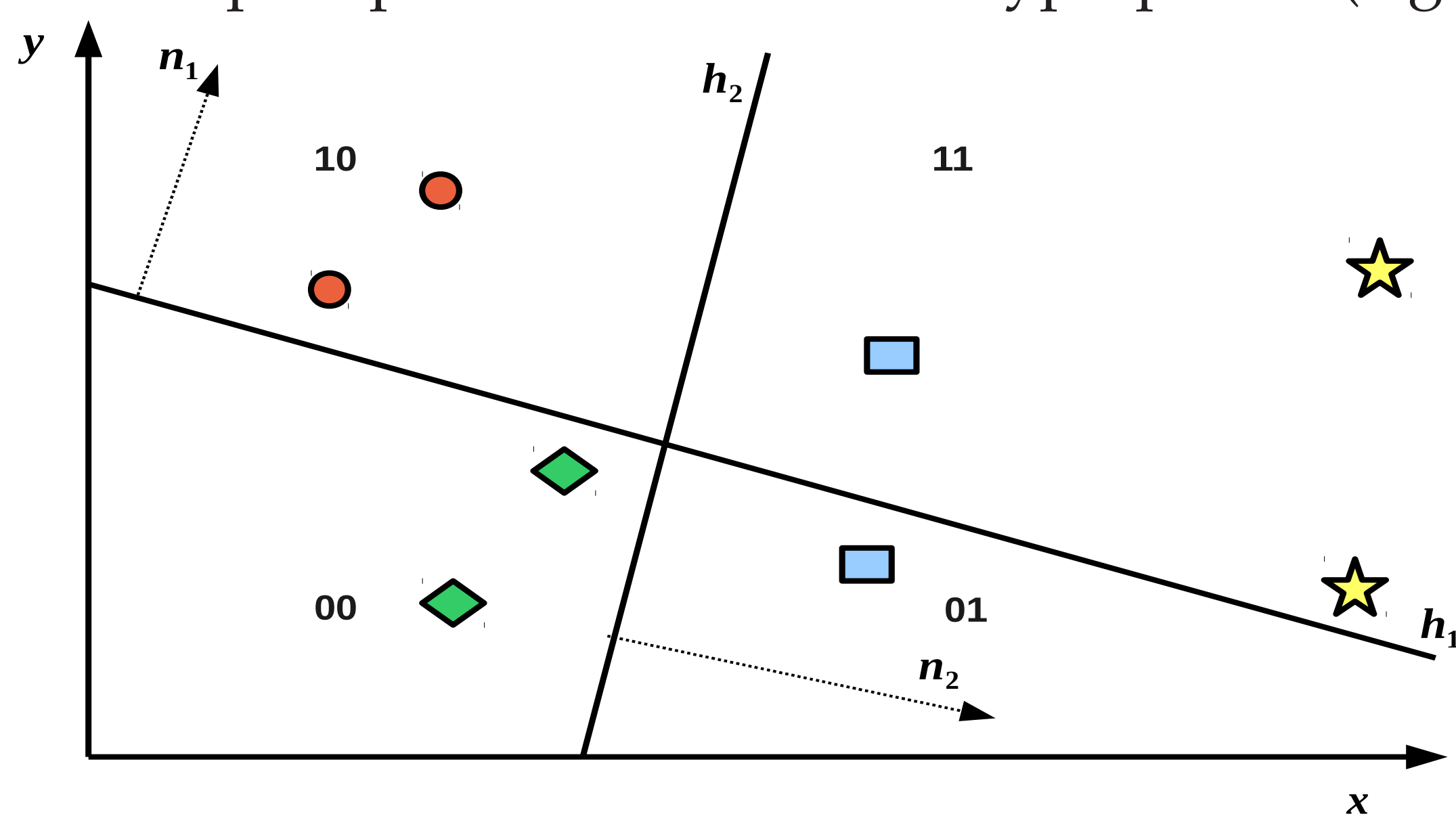
- Problem:** Fast Nearest Neighbour (NN) search in large datasets.
- Hashing-based approach:**
 - Generate a similarity preserving binary code (fingerprint).
 - Use fingerprint as index into the buckets of a hash table.
 - If collision occurs only compare to items in the same bucket.



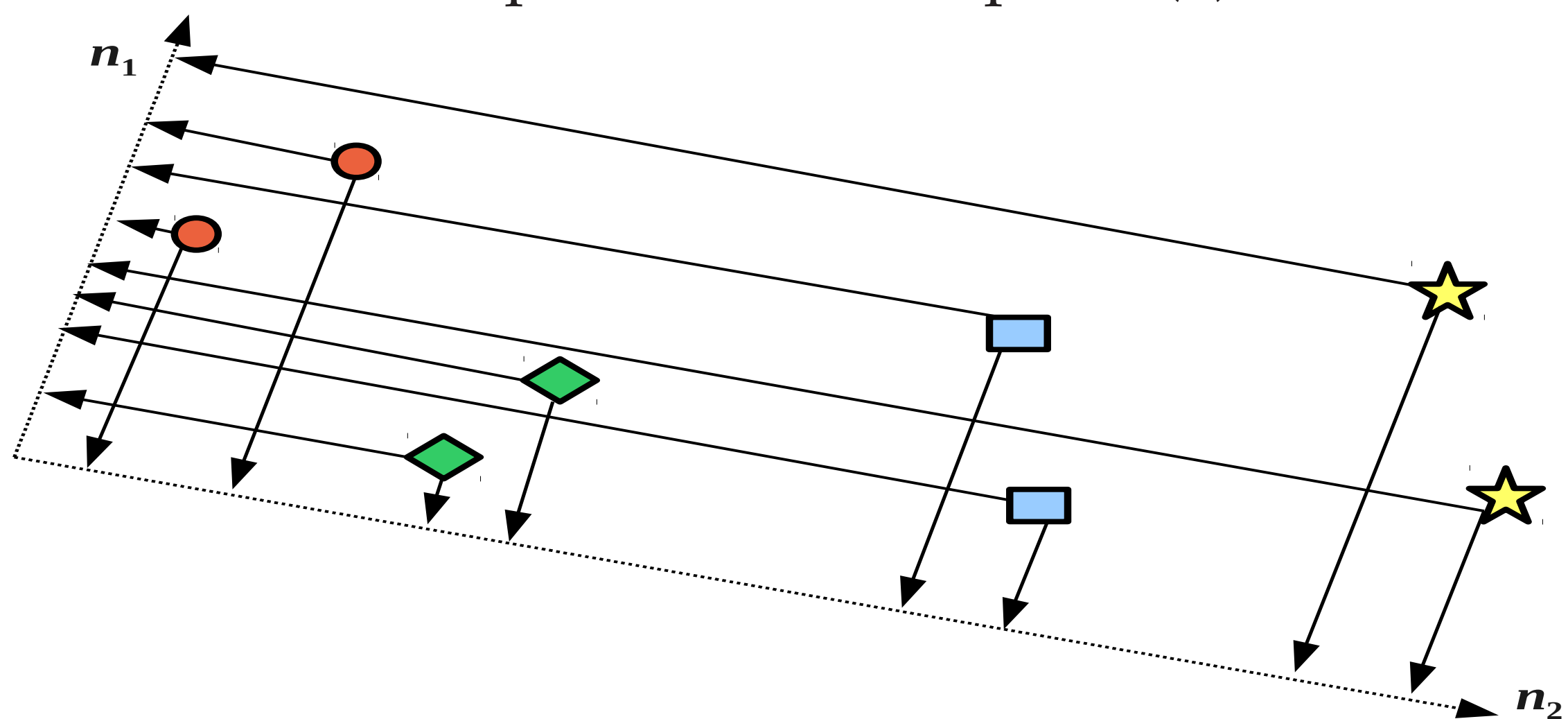
- Advantages:**
 - Constant query time (with respect to the database size).
 - Compact binary codes are extremely storage efficient.

LOCALITY SENSITIVE HASHING (LSH) (INDYK AND MOTWANI, '98)

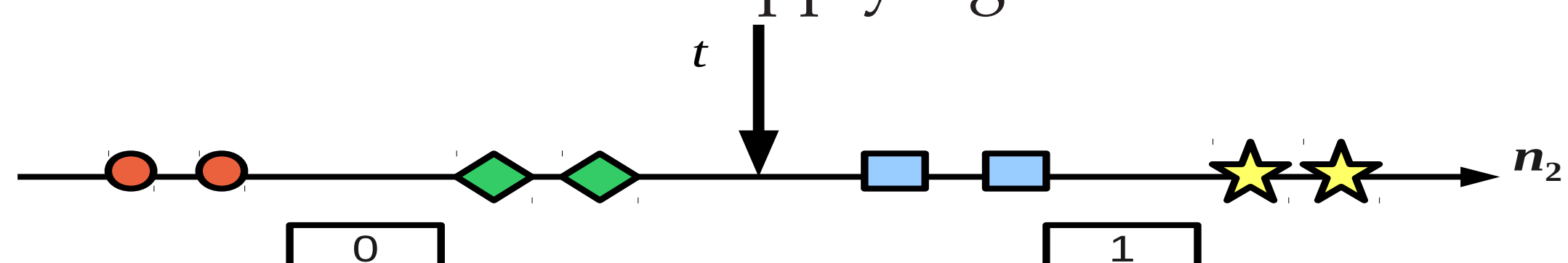
- Randomized algorithm for approximate nearest neighbour (ANN) search using binary codes.
- Probabilistic guarantee on retrieval accuracy versus search time.
- LSH for inner product similarity:
 - Divide input space with L random hyperplanes (e.g. L=2):



- Projection:** Take dot product of data-point (x) with normal ($n.x$):



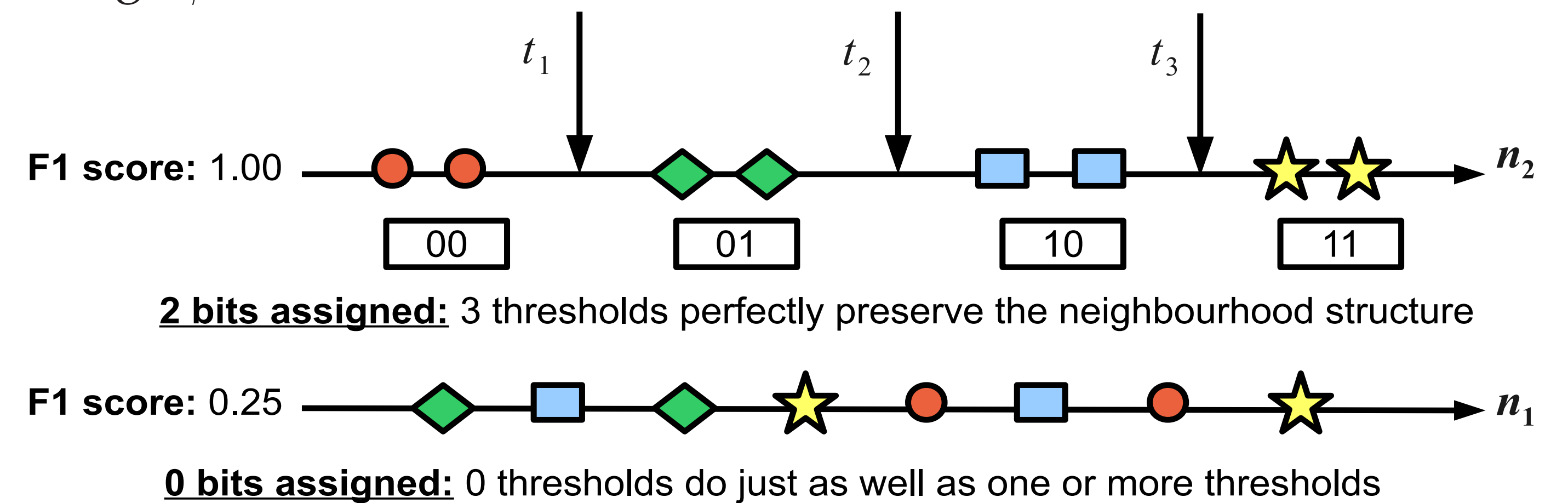
- Quantisation:** Threshold and apply sign function:



- Vanilla LSH:** each hyperplane gives 1 bit of the binary code.

VARIABLE BIT QUANTISATION FOR LSH

- Goal: assign more bits to neighbourhood preserving hyperplanes.
- Score hyperplanes by F_β computed using pairwise constraints matrix S : if $S_{ij} = 1$ then points x_i, x_j with projections y_i, y_j are true nearest neighbours.
- TP: $\# S_{ij} = 1$ pairs in same region. FP: $\# S_{ij} = 0$ pairs in same region. FN: $\# S_{ij} = 1$ pairs in different regions. Combine TP, FP, FN using F_β :



- Solve bit allocation as a binary integer linear program (BILP):

$$\begin{aligned} \max \quad & \|F \circ Z\| \\ \text{subject to} \quad & \|Z_h\| = 1 \quad h \in \{1 \dots B\} \\ & \|Z \circ D\| \leq B \\ & Z \text{ is binary} \end{aligned}$$

$$\begin{matrix} F & h_1 & h_2 & D & Z \\ b_0 & \begin{pmatrix} 0.25 & 0.25 \end{pmatrix} & \begin{pmatrix} 0 & 0 \end{pmatrix} & \begin{pmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 1 \end{pmatrix} \\ b_1 & \begin{pmatrix} 0.35 & 0.50 \end{pmatrix} & \begin{pmatrix} 1 & 1 \end{pmatrix} \\ b_2 & \begin{pmatrix} 0.40 & 1.00 \end{pmatrix} & \begin{pmatrix} 2 & 2 \end{pmatrix} \end{matrix}$$

- F : matrix of F_β scores. Z : bit allocation. D : constraint matrix. B : bit budget. $\|\cdot\|$: Frobenius L_1 norm. $\circ$: Hadamard product

EVALUATION PROTOCOL

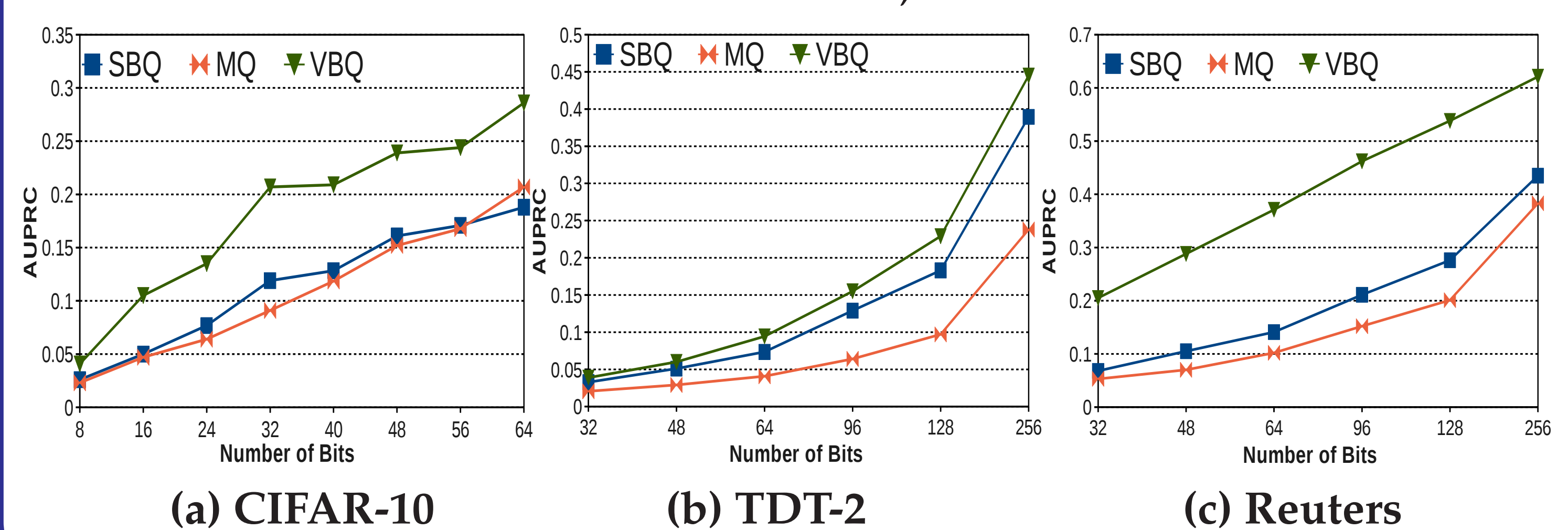
- Task:** Retrieval on: *CIFAR-10*, *Reuters-21578* and *TDT-2* datasets.
- Projections:** LSH, Shift-invariant kernel hashing (SIKH), Binary-LSI (BLSI), Spectral Hashing (SH) and PCA-Hashing (PCAH).
- Baselines:** Single Bit Quantisation (SBQ), Manhattan Hashing (MQ) (Kong et al., '12), Double-Bit quantisation (DBQ) (Kong and Li, '12).
- Hamming Ranking:** how well do we retrieve ϵ -NN of queries? Quantify using area under the precision-recall curve (AUPRC).

RESULTS

- AUPRC for projections at 32 bits (images) and 128 bits (text):

Dataset	CIFAR-10				TDT-2				Reuters			
	SBQ	MQ	DBQ	VBQ	SBQ	MQ	DBQ	VBQ	SBQ	MQ	DBQ	VBQ
SIKH	0.042	0.063	0.047	0.161	0.034	0.045	0.031	0.092	0.102	0.112	0.087	0.389
LSH	0.119	0.093	0.066	0.207	0.189	0.097	0.089	0.229	0.276	0.201	0.175	0.538
BLSI	0.038	0.135	0.111	0.231	0.283	0.210	0.087	0.396	0.100	0.030	0.030	0.156
SH	0.051	0.135	0.111	0.202	0.146	0.212	0.167	0.370	0.033	0.028	0.030	0.154
PCAH	0.036	0.137	0.107	0.219	0.281	0.208	0.094	0.374	0.095	0.034	0.027	0.154

- AUPRC vs. number of bits for CIFAR, TDT-2 and Reuters:



FUTURE WORK

- Evaluation of VBQ in a hash lookup based retrieval scenario.